

Mobile Communications

TCS 455

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Lecture 11

Office Hours:

BKD 3601-7

Tuesday 14:00-16:00

Thursday 9:30-11:30

Announcements

- Read
 - Chapter 3: 3.1 – 3.2, 3.5.1, 3.6, 3.7.2
 - Posted on the web
 - Appendix A.1 (Erlang B)

Big Picture

S = total # available duplex radio channels for the system



Frequency reuse with **cluster size N**

“Capacity”

$$C = \frac{A_{\text{total}}}{A_{\text{cell}}} \times \frac{S}{N}$$

Tradeoff

$$\frac{S}{I} \approx \frac{kR^{-\gamma}}{K \times (kD^{-\gamma})} = \frac{1}{K} \left(\frac{D}{R} \right)^\gamma = \frac{1}{K} (\sqrt{3N})^\gamma$$

m = # channels allocated to each cell.

- Omni-directional: $K = 6$
- 120° Sectoring: $K = 2$
- 60° Sectoring: $K = 1$



Trunking

$$P_b = \frac{\frac{A^m}{m!}}{\sum_{i=0}^m \frac{A^i}{i!}}$$

λ = Average # call attempts/requests per unit time

A = **traffic intensity** or load [Erlangs] = $\frac{\lambda}{\mu}$

$\frac{1}{\mu} = H$ = Average call length

Call blocking probability

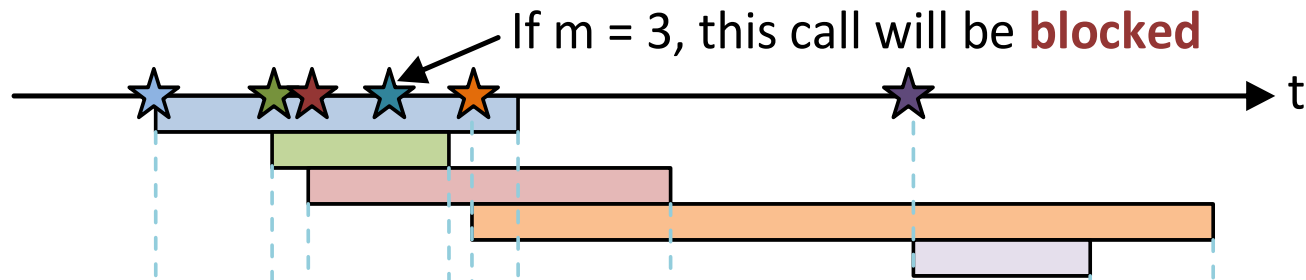
Erlang-B formula

Assumption

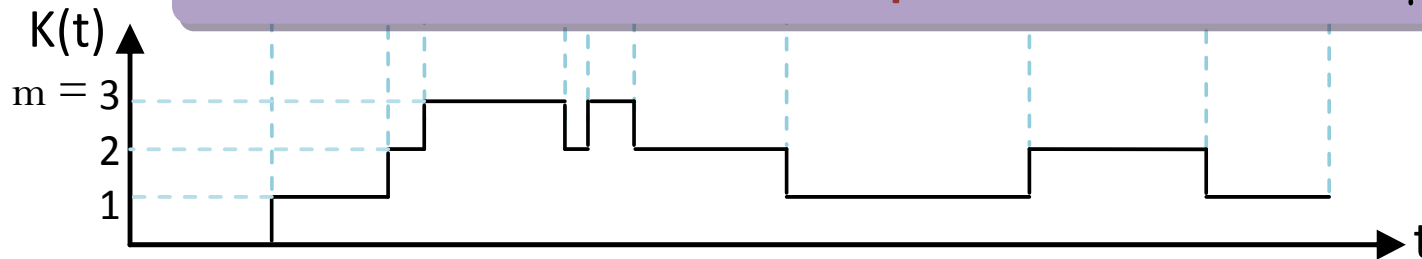
- **Blocked calls cleared**
 - Offers no queuing for call requests.
 - For every user who requests service, it is assumed there is no setup time and the user is given immediate access to a channel if one is available.
 - If no channels are available, the requesting user is blocked without access and is free to try again later.
- **Calls arrive as determined by a *Poisson process*.**
- There are memoryless arrivals of requests, implying that all users, including blocked users, may request a channel at any time.
- There are an infinite number of users (with finite overall request rate).
 - The finite user results always predict a smaller likelihood of blocking. So, assuming infinite number of users provides a conservative estimate.
- **The duration of the time that a user occupies a channel is exponentially distributed**, so that longer calls are less likely to occur.
- There are m channels available in the trunking pool.
 - For us, $m =$ the number of channels for a cell (C) or for a sector

Assumption (2)

The call request process is **Poisson** with rate λ



The duration of calls are i.i.d. **exponential** r.v. with rate μ .

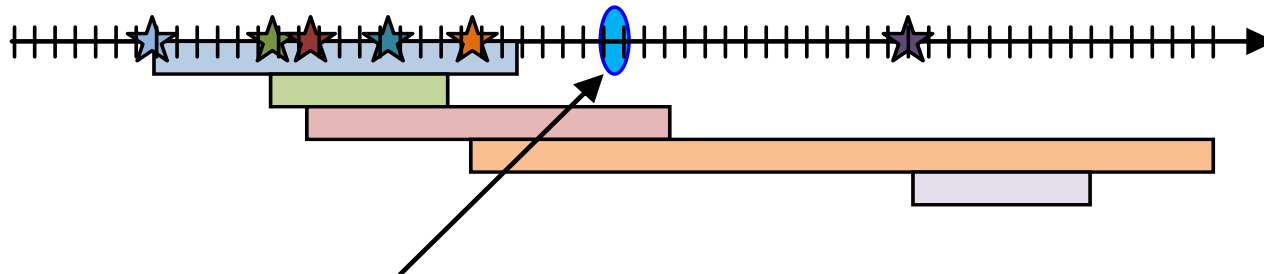


$K(t)$ = "state" of the system
= the number of used channel at time t

We want to find out what proportion of time the system has $K = m$.

Small Slot Analysis

Suppose each slot duration is δ .



- Consider the i^{th} small slot.
- Let $K_i = k$ be the value of K at the beginning of this time slot.
- $k = 2$ in the above figure.
- Then, K_{i+1} is the value of K at the end of this slot which is the same as the value of K at the beginning of the next slot.

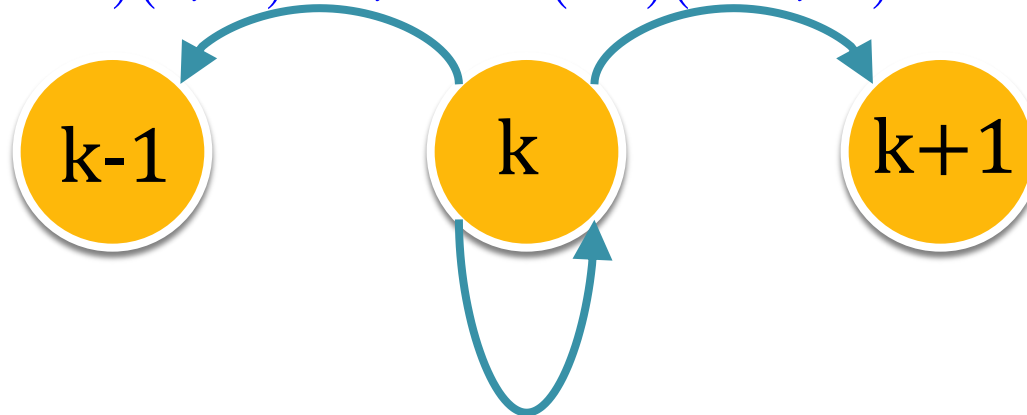
- $P[0 \text{ new call request}] \approx 1 - \lambda\delta$
- $P[1 \text{ new call request}] \approx \lambda\delta$
- $P[0 \text{ old-call end}] \approx (1 - \mu\delta)^k \approx 1 - k\mu\delta$
- $P[1 \text{ old-call end}] \approx k\mu\delta(1 - \mu\delta)^{k-1} \approx k\mu\delta$

How do these events affect K_{i+1} ?

Small slot Analysis (2)

$$K_{i+1} = K_i + (\# \text{ new call request}) - (\# \text{ old-call end})$$

$$(1 - \lambda\delta)(k\mu\delta) \approx k\mu\delta \quad (\lambda\delta)(1 - k\mu\delta) \approx \lambda\delta$$



$$(1 - \lambda\delta)(1 - k\mu\delta) + (\lambda\delta)(k\mu\delta) \approx 1 - \lambda\delta - k\mu\delta$$

The labels on the arrows are probabilities.

$$P[0 \text{ new call request}] \approx 1 - \lambda\delta$$

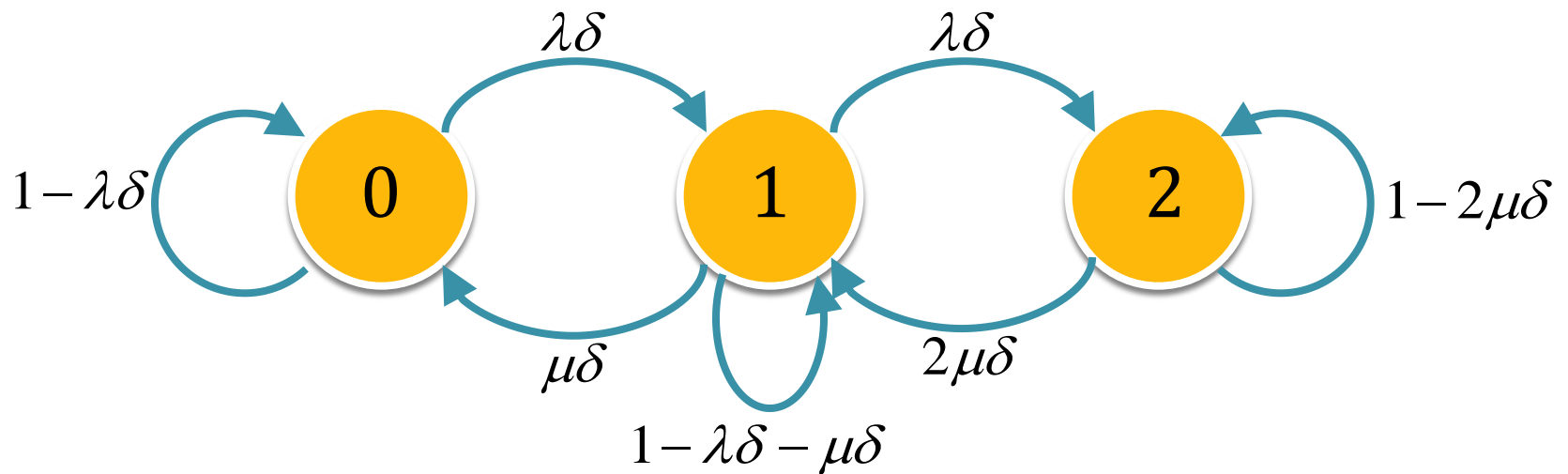
$$P[1 \text{ new call request}] \approx \lambda\delta$$

$$P[0 \text{ old-call end}] \approx 1 - k\mu\delta$$

$$P[1 \text{ old-call end}] \approx k\mu\delta$$

Small slot Analysis: Markov Chain

- Case: $m = 2$



Markov Chain

- Markov chains model many phenomena of interest.
- We will see one important property: **Memoryless**
 - It retains no memory of where it has been in the past.
 - Only the current state of the process can influence where it goes next.
- Very similar to the *state transition diagram* in digital circuits.
 - In digital circuit, the labels on the arrows indicate the input/control signal.
 - Here, the labels on the arrows indicate transition probabilities. (If the system is currently at a particular state, where would it go next on the next time slot?)
- We will focus on **discrete time Markov chain**.

Example: The Land of Oz

- Land of Oz is blessed by many things, but not by good weather.
 - They never have two nice days in a row.
 - If they have a nice day, they are just as likely to have snow as rain the next day.
 - If they have snow or rain, they have an even chance of having the same the next day.
 - If there is change from snow or rain, only half of the time is this a change to a nice day.
- If you visit the land of Oz next year for one day, what is the chance that it will be a nice day?

The Land of Oz: Transition Matrix

$$P = \begin{matrix} & \begin{matrix} \text{R} & \text{N} & \text{S} \end{matrix} \\ \begin{matrix} \text{R} \\ \text{N} \\ \text{S} \end{matrix} & \begin{bmatrix} \frac{1}{2} & \frac{1}{4} & \frac{1}{4} \\ \frac{1}{2} & 0 & \frac{1}{2} \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{2} \end{bmatrix} \end{matrix}$$

$$P[K_{i+1} = \text{rain} | K_i = \text{normal}]$$